

Fast Compressive Phase Retrieval under Bounded Noise

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Proof of Measure of Concentration

Lemma 1 ((Shalev-Shwartz and Ben-David 2014)). *Let \mathcal{F} be the class of linear predictors with the L_2 norm of the weights bounded by W_2 . Assume that the L_2 norm of the instance is bounded by X_2 . Then for the ρ -Lipschitz loss ℓ such that $\max_{(w,x) \in [-W_2 X_2, W_2 X_2]} |\ell(w, x, y)| < U$, with probability at least $1 - \delta$ over the choice of an i.i.d. sample T of size m ,*

$$\forall w \in \{w : \|w\|_2 \leq W_2\},$$

$$|\mathbb{E}\ell(w, x, y) - \ell(w, T)| \leq \frac{2\rho W_2 X_2}{\sqrt{m}} + U \sqrt{\frac{2 \log(2/\delta)}{m}}.$$

Lemma 2 (Lemma 3 in Main Body). *Let $\mathbf{z} \in \{\mathbf{z} : \|\mathbf{z}\|_2 \leq 1\}$ and $\{\mathbf{w}_i\}_{i=1}^m$ be random vectors i.i.d. sampled from the standard Gaussian distribution $\mathcal{N}(\mathbf{0}, \mathbf{I})$. Fix $\mathbf{x} \in \mathbb{R}^n$ and suppose that $m \geq c_0 d \log^4(\frac{1}{\delta}) \epsilon^{-2}$ with a universal constant c_0 , then with probability at least $1 - \delta$,*

$$|f_{\mathbf{x}_0}(\mathbf{z}) - \mathbb{E}f_{\mathbf{x}_0}(\mathbf{z})| \leq \epsilon \quad (1)$$

uniformly holds for all $\mathbf{z} \in \mathbb{R}^n$.

Proof. The proof is basically based on Lemma 1. Note that

$$f_{\mathbf{x}_0}(\mathbf{z}) = \frac{1}{m} \sum_{i=1}^m y_i (\mathbf{w}_i^T \mathbf{z})^2 = \frac{1}{m} \sum_{i=1}^m ((\mathbf{w}_i^T \Psi \mathbf{x}_0)^2 + \eta_i) (\mathbf{w}_i^T \mathbf{z})^2. \quad (2)$$

By Lemma 1, we have that

$$\Pr \left[\sup_{\mathbf{z}} |\mathbb{E}f_{\mathbf{x}}(\mathbf{z}) - f_{\mathbf{x}}(\mathbf{z})| > \frac{2\rho W_2 X_2}{\sqrt{m}} + s \right] \leq 2 \exp \left(-\frac{ms^2}{2U^2} \right), \quad (3)$$

where the supremum is taken over all $\mathbf{z} \in \{\mathbf{z} : \|\mathbf{z}\|_2 \leq 1\}$.

To identify the parameters ρ , W_2 , X_2 , and U above, we exploit the property of standard Gaussian distribution. Specifically, we see that $W_2 = 1$. By Lemma 4, we have $\|\mathbf{w}_i\|_2 \leq O(\sqrt{d \log(1/\delta)}) \triangleq X_2$ with probability at

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least $1 - \delta$. Let $\ell(\mathbf{w}^T \mathbf{z}) = y_i (\mathbf{w}^T \mathbf{z})^2$. Since by Lemma 3 and the fact that η_i is a global constant, $|\ell'(\mathbf{w}^T \mathbf{z})| \leq 2y_i |\mathbf{w}^T \mathbf{z}| \leq O(\log^{3/2}(1/\delta))$ for any $\mathbf{w}^T \mathbf{z}$ with high probability. So $\ell(\mathbf{w}^T \mathbf{z})$ is $O(\log^{3/2}(1/\delta))$ -Lipschitz, i.e., $\rho = O(\log^{3/2}(1/\delta))$. Furthermore, $|\ell(\mathbf{a}^T \mathbf{z})| \leq O(\log^2(1/\delta)) \triangleq U$. Plugging in all those parameters, we can see that when $m \geq c_0 d \log^4(\frac{1}{\delta}) \epsilon^{-2}$, the R.H.S. of (3) is no larger than δ , as desired. \square

Property of Standard Gaussian Distribution

Lemma 3. *Let X be the random variable drawn from standard Gaussian distribution $\mathcal{N}(0, 1)$. Then for every $t > 0$,*

$$\Pr[|X| > t] \leq \exp(-t^2/2). \quad (4)$$

Lemma 4. *Let \mathcal{P} be the isotropic Gaussian distribution in \mathbb{R}^d . Then $\Pr_{\mathbf{w} \sim \mathcal{P}}[\|\mathbf{w}\|_2 \geq \alpha] \leq \left(\frac{e\alpha^2}{d}\right)^{d/2} e^{-\alpha^2/2}$.*

Proof. We have

$$\begin{aligned} \Pr[\|\mathbf{w}\|_2 \geq \alpha] &= \Pr[e^{s\|\mathbf{w}\|_2^2} \geq e^{s\alpha^2}] \\ &\leq \frac{\mathbb{E}e^{s\|\mathbf{w}\|_2^2}}{e^{s\alpha^2}} \\ &= e^{-s\alpha^2} (1 - 2s)^{-d/2}, \end{aligned} \quad (5)$$

where the last equality is from the moment generating function of Chi-Square distribution. Setting $s = \frac{\alpha^2 - d}{2\alpha^2}$, we obtain the desired result. \square

Result on Standard Compressed Sensing

Theorem 5 ((Foucart and Rauhut 2013), Robust Recovery). *Let $\mathbf{x} \in \mathbb{R}^n$ be a t -sparse vector. Suppose that $\mathbf{A} \in \mathbb{R}^{m \times n}$ is a randomly drawn standard Gaussian matrix. Assume that the noisy measurements $\mathbf{b} = \mathbf{A}\mathbf{x} + \mathbf{e}$ are taken with $\|\mathbf{e}\|_2 \leq \sqrt{\eta/m}$. If*

$$\frac{m^2}{m+1} \geq 2t \left(\sqrt{\log(en/t)} + \sqrt{\log(\delta^{-1})/t} + \tau/\sqrt{t} \right)^2, \quad (6)$$

then with probability at least $1 - \delta$, every minimizer $\tilde{\mathbf{x}}$ to

$$\tilde{\mathbf{x}} := \min_{\mathbf{x}} \|\mathbf{x}\|_1, \quad \text{s.t.} \quad \|\tilde{\mathbf{b}} - \Psi \mathbf{x}\|_2 \leq \sqrt{\eta/m}. \quad (7)$$

satisfies

$$\|\tilde{\mathbf{x}} - \mathbf{x}\|_2 \leq 2 \frac{\sqrt{\eta}}{\tau\sqrt{m}}. \quad (8)$$

Property of Bernoulli Model

Lemma 6. *Let n be the number of Bernoulli trials and suppose that $\Omega \sim \text{Ber}(d/n)$. Then with probability at least $1 - \delta$, $|\Omega| = \Theta(d)$, provided that $d \geq 4 \log(1/\delta)$.*

Proof. Take a perturbation ϵ such that $d/n = d_0/n + \epsilon$. By the scalar Chernoff bound which states that

$$\Pr(|\Omega| \leq d_0) \leq e^{-\epsilon^2 n^2 / 2d_0}, \quad (9)$$

if taking $d_0 = d/2$, $\epsilon = d/2n$ and $d \geq 4 \log(1/\delta)$, we have

$$\Pr(|\Omega| \leq d/2) \leq e^{-d/4} \leq \delta. \quad (10)$$

On the other hand, by the scalar Chernoff bound again which states that

$$\Pr(|\Omega| \geq d_0) \leq e^{-\epsilon^2 n^2 / 3d}, \quad (11)$$

if taking $d_0 = 2d$, $\epsilon = -d/n$ and $d \geq 4 \log(1/\delta)$, we obtain

$$\Pr(|\Omega| \geq 2d) \leq e^{-d/3} \leq \delta. \quad (12)$$

Finally, according to (10) and (12), we conclude that $d/2 < |\Omega| < 2d$ with probability at least $1 - \delta$. \square

References

- Foucart, S., and Rauhut, H. 2013. *A Mathematical Introduction to Compressive Sensing*, volume 1. Springer.
- Shalev-Shwartz, S., and Ben-David, S. 2014. *Understanding Machine Learning: From Theory to Algorithms*. Cambridge University Press.